The Expressiveness of Session Types

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It is fair to say that expressiveness has been key to the success and impact of session types: enhancing the expressivity of typed processes and/or the properties enforced by typing is often a strong motivation for developing new typed frameworks. Rigorously contrasting different variants of session types from an expressiveness perspective is a pressing and non-trivial challenge. This short note gives a unified overview of some of our works in this direction; they cover various angles, including: higher-order communication, binary and multiparty protocols, and “propositions-as-sessions”. The talk at ST30 will reflect on these (recent) results but also on avenues for future exploration.

In a 2016 position paper [18], I argued about the challenge of typed expressiveness in concurrency. The challenge consists in lifting the techniques that have proved so successful for assessing the expressiveness of untyped π-calculi to the setting of (session) typed π-calculi. The motivations are similar as in the untyped case: we would like to understand the relative strengths of different type systems, and to compare them rigorously to understand their differences and relative merits. This is actually a dynamic problem, as the number of behavioral type systems appearing in the literature is ever increasing. By disciplining the use of channels/names in processes, (session) types strongly influence their expressiveness. Casting expressiveness studies in the typed setting puts the focus on the correctness properties ensured by typing: different type systems enforce different properties, and so they determine a new, crucial dimension for comparisons. This is relevant in order to have satisfactory ‘apples-to-apples’ comparisons.

We have pursued the challenge of typed expressiveness particularly within the project “Unifying Correctness for Communicating Software” (2019-2024), whose goal is to connect different type systems via new expressiveness results. The talk at ST30 will give a unified overview of some selected results:

**Higher-Order Concurrency** Session π-calculi with higher-order communication, in which exchanged values include processes, are interesting because they can specify code mobility. In [13,15], we study a family of languages with higher-order communication. We identify a core higher-order session calculus, which supports only passing of abstractions (functions from names to processes) but lacks recursion. Still, it can correctly codify name passing, recursion, and higher-order abstractions (functions from processes to processes). We obtain a taxonomy of session-typed languages, defined by a series of typed encodability results: each result consists of a translation on processes (as usual) but also of a translation on types, which ensures that source protocols are properly translated. In all cases, types are essential in defining the translations and/or in proving their correctness. Defining this taxonomy requires (i) enhancing encodability criteria with session types and (ii) developing (typed) behavioral equivalences [14].

**Minimal Session Types** The study in [13,15] considers a standard formulation of session types. In [3,2], we explore a simpler formulation, which dispenses with sequencing at the level of types. This way, e.g., in a type such as ‘!⟨S⟩; T’ the continuation type T can only be ‘end’. This formulation, dubbed minimal session types, resembles simple and linear types for the π-calculus but also the way in which (one-shot) channels are typically used in languages such as Go. Using the core higher-order session calculus from [13,15], we prove that sequencing constructs in processes and session types is convenient but
redundant: every process typable with standard session types can be compiled down to a process typable with minimal session types. This shows that only sequenti\text{ality} in processes is truly indispensable, as it can correctly codify sequenti\text{ality} in types. This can be considered as a result of absolute expressiveness, which explains session types in terms of themselves, without appealing to extraneous type disciplines. The idea behind minimal session types is rather simple and yet robust, as the formulation and associated correctness results can be adapted also to sessions the first-order setting, as shown in [1].

**Binary and Multiparty Session Types** Arguably, the two most active research strands within session types concern multiparty session types and “propositions-as-sessions”—the propositions-as-types approach to session types. The work in [4] jointly addresses them by developing an analysis of multiparty protocols using the binary session type system derived from the Curry-Howard interpretation of linear logic as session types. The idea is to analyze a multiparty implementation (a collection of processes, one per protocol participant) together with a so-called medium process, which is synthesized from a given global type. Although our analysis is not an encodability result as usual, it does bear witness to the expressivity of the session types based on linear logic, and enables to transfer the deadlock-freedom for processes involving delegation and interleaving to the multiparty setting. The analysis has been generalized in [10], which develops a technique based on router processes, which “wrap” local implementations while enabling the composition of protocol participants in arbitrary topologies. Routers can be used in static verification (as shown in [10]) but also in run-time verification, as recently shown in [11].

**Type Systems for Deadlock-Freedom** The work [7, 8] compares two type systems that enforce the deadlock-freedom property: (i) a type system based on priorities, as pioneered by Kobayashi [12]; and (ii) a type system based on propositions-as-sessions. To enable a coherent, ‘apples-to-apples’ comparison, two classes of processes are defined, denoted $\mathcal{K}$ and $\mathcal{L}$, which contain the deadlock-free induced by the type systems of [12] and [19], respectively. One key result in [7, 8] is that $\mathcal{L} \subseteq \mathcal{K}$, i.e., the priority-based approach to deadlock-freedom subsumes the proof-theoretical approach. A key insight is that while the class $\mathcal{L}$ contains exclusively processes with tree-like topologies, the class $\mathcal{K}$ contains also processes with safe forms of circular topologies. This work also gives translations of $\mathcal{K}$ into $\mathcal{L}$.

**Typed Functions into Session-Typed Processes** Encodings of the $\lambda$-calculus provide a significant stress test of the expressiveness of the $\pi$-calculus. Following this line, we have developed encodings of resource $\lambda$-caluli into session-typed $\pi$-caluli derived from propositions-as-sessions [16, 17]. A distinctive aspect here is non-determinism: in $\lambda$, non-determinism concerns fetching of resources from bags of available resources; in $\pi$, we have sessions that are non-deterministically available—they may be available but may also fail. Because our resource $\lambda$-caluli are equipped with intersection types, our encodability results delineate a new, surprising connection between intersection types and linear logic.

**Beyond Linear Logic** The discovery of ‘propositions-as-sessions’ has a substantial and direct impact on our research program on typed expressiveness, for two reasons. First, as the correspondence naturally entails several key correctness properties (fidelity, deadlock-freedom, confluence, strong normalization), it defines a canonical class of session processes, upon which all other classes should be compared. Second, logic suggests a principled avenue for extensions: new typed frameworks can be obtained by suitably extending the underlying (linear) logic. Representative examples of this kind of extensions include behavioral polymorphism [5] and domain-aware computation [6]; just right outside linear logic, recent work has established a new concurrent interpretation of the logic of bunched implications [9].

References


